

NEIGHBORHOOD HYPERGRAPHS OF DIGRAPHS
AND SOME MATRIX PERMUTATION PROBLEMS

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NEIGHBORHOOD HYPERGRAPHS OF DIGRAPHS AND SOME MATRIX PERMUTATION PROBLEMS

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Abstract. Given a digraph D , the set of all pairs $(N^-(v), N^+(v))$ constitutes the neighborhood dihypergraph $\mathcal{N}(D)$ of D . The Digraph Realization Problem asks whether a given dihypergraph H coincides with $\mathcal{N}(D)$ for some digraph D . This problem was introduced by Aigner and Triesch [2] as a natural generalization of the Open Neighborhood Realization Problem for undirected graphs, which is known to be NP-complete.

We show that the Digraph Realization Problem remains NP-complete for orgraphs (orientations of undirected graphs). As a corollary, we show that the Matrix Skew-Symmetrization Problem for square $\{0, 1, -1\}$ matrices $(a_{ij} = -a_{ji})$ is NP-complete. This result can be compared with the known fact that the Matrix Symmetrization Problem for square $0 - 1$ matrices $(a_{ij} = a_{ji})$ is NP-complete.

Extending a negative result of Fomin, Kratochvíl, Lokshtanov, Mancini, and Telle [15] we show that the Digraph Realization Problem remains NP-complete for almost all hereditary classes of digraphs defined by a unique minimal forbidden subdigraph.

Finally, we consider the Matrix Complementation Problem for rectangular $0 - 1$ matrices, and prove that it is polynomial-time equivalent to graph isomorphism. A related known result is that the Matrix Transposability Problem is polynomial-time equivalent to graph isomorphism.

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05C85 (Graph algorithms).

Neighborhood hypergraphs of digraphs and orgraphs, Graph Isomorphism Problem, matrix symmetrization, matrix complementation, symmetrizability, skew-symmetrizability, involutory automorphisms

1 Introduction

Let $D = (V, A)$ be a digraph without loops and multiple arcs. For a vertex $v \in V$, we denote

$$N^-(v) = \{u \in V : (u, v) \in A\},$$

the *in-neighborhood* of v , and

$$N^+(v) = \{w \in V : (v, w) \in A\},$$

the *out-neighborhood* of v . Suppose that we know all pairs $(N^-(v), N^+(v))$, is it possible to restore the digraph? To formalize the problem, let us define a *directed hypergraph*, or shortly *dihypergraph*, as an ordered pair $(V, A) = H$ consisting of a finite set V , the *vertex-set* of H , and a finite multi-set of hyperarcs, a *hyperarc* $a \in A$ being an ordered pair $(a^-, a^+) = a$ of some subsets a^- and a^+ of V . It is possible that $a^- = \emptyset$ or $a^+ = \emptyset$ or $a^- = a^+$. Also note that a^- and a^+ are not necessarily disjoint.

Definition 1 *The neighborhood dihypergraph of a digraph $D = (V, A)$, $\mathcal{N}(D)$, has V as its vertex-set, and $A(\mathcal{N}(D)) = \{(N^-(v), N^+(v)) : v \in V\}$.*

An obvious property of $\mathcal{N}(D)$ is that the number of vertices is the same as the number of hyperarcs. The following problem was proposed by Aigner and Triesch [2].

Decision Problem 1 (Digraph Realization Problem)

Instance: A directed hypergraph H .

Question: Does $H = \mathcal{N}(D)$ hold for some digraph D ?

This problem generalizes the *Open Neighborhood Realization Problem* for undirected graphs: given a hypergraph H (with possible multiple hyperedges), the problem is asking to find a graph G for which H is the hypergraph of open neighborhoods $\mathcal{N}^{\text{op}}(G)$, of vertices of G , that is $V(H) = V(G)$ and $E(H) = \{N(v) : v \in V(G)\}$. Here $N(v) = \{w \in V(G) : vw \in E(G)\}$ is the *neighborhood* of a vertex v of G . The Open Neighborhood Realization Problem was proposed by Sós [21] under the name the Star System Problem, and it is also attributed to G. Sabidussi by Babai [4]. Also, Babai [4] noticed that the problem is at least as hard as graph isomorphism. The *Graph Isomorphism Problem* is well-known: Are two given graphs isomorphic? Boros, Gurvich, and Zverovich [8] survey different equivalent formulations of the problem.

The *Closed Neighborhood Realization Problem* is defined in a similar way, using the *closed neighborhoods* $N[v] = \{v\} \cup N(v)$ of vertices. Also, one can consider a hypergraph $\mathcal{N}(G)$ of open and closed neighborhoods of G , that is, for each vertex v either $N(v)$ or $N[v]$ is a hyperedge of $\mathcal{N}(G)$. The *Neighborhood Realization Problem* is to decide whether a given hypergraph H is $\mathcal{N}(G)$ for some graph G .

Theorem 1 (Lalonde [16, 17]) *The Open Neighborhood Realization Problem, the Closed Neighborhood Realization Problem, and the Neighborhood Realization Problem are NP-complete.*

An undirected graph G can be viewed as a digraph on $V(G)$ if we replace every edge $uv \in E(G)$ by the corresponding pair (u, v) , (v, u) of opposite arcs.

Corollary 1 (Aigner and Triesch [2]) *The Digraph Realization Problem is NP-complete.*

Theorem 1 has an interesting interpretation. A square matrix $A = (a_{ij})$ is *symmetric* if $a_{ij} = a_{ji}$ for all i and j . A square matrix A is *symmetrizable* if it is possible to permute rows of A in such a way that the resulting matrix is symmetric. The Neighborhood Realization Problem is equivalent to the *Matrix*

Symmetrization Problem: Is a given square 0 – 1 matrix is symmetrizable? If we additionally require that all entries in the main diagonal are 0s (respectively, 1s), then we obtain a problem which is equivalent to the Open (respectively, Closed) Neighborhood Realization Problem. The three symmetrization problems are NP-complete.

We show that the Digraph Realization Problem remains NP-complete for orgraphs (orientations of undirected graphs) and for almost all hereditary classes of digraphs defined by a unique minimal forbidden subdigraph. As a corollary, we show that the Matrix Skew-Symmetrization Problem for square $\{0, 1, -1\}$ matrices is NP-complete. The problem is to bring a matrix to skew form ($a_{ij} = -a_{ji}$) using permutations of rows. Then we consider the Matrix Complementation Problem for rectangular 0 – 1 matrices: to construct the complementary matrix (defined by $\bar{a}_{ij} = 1 - a_{ij}$) using row and column permutations. We prove that it is polynomial-time equivalent to graph isomorphism.

2 Representations

It is convenient to represent hypergraphs as bipartite graphs. and as their adjacency matrices. A *bigraph* $B = (X, Y, E)$ is defined as a bipartite graph on vertex-set $V = X \cup Y$ with a fixed order (X, Y) of its parts. Here $X \cap Y = \emptyset$ and $E \subseteq X \times Y$. To a bigraph $B = (X, Y, E)$ we can associate its *X-Y-adjacency matrix* $A(B) = (a_{ij}) \in \{0, 1\}^{X \times Y}$ defined by $a_{ij} = 1$ if and only if $(i, j) \in E$. Conversely, any 0 – 1 matrix $A = (a_{ij})$ can be viewed as the X-Y adjacency matrix $A = A(B)$ of a corresponding bigraph $B = (X, Y, E)$, where X is the set of row indices of A , Y is the set of column indices of A , and $(i, j) \in E$ if and only if $a_{ij} = 1$, see an example in Figure 1.

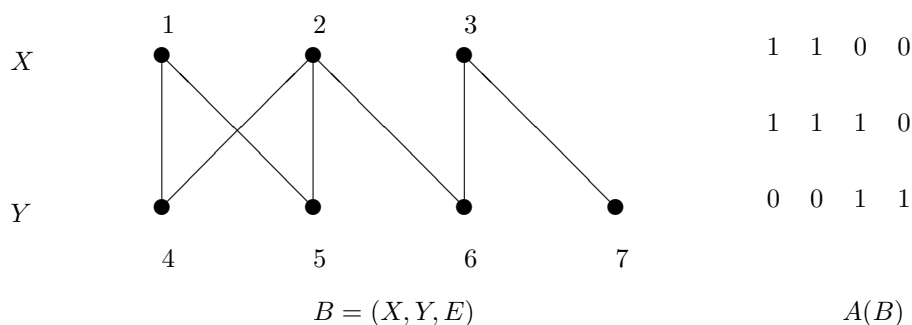


Figure 1: A bigraph $B = (X, Y, E)$ and its adjacency matrix $A(B)$.

Now we consider similar representations of a dihypergraph H . Let us define a *directed bigraph* $B = (X, Y, A)$ as a bipartite digraph on vertex-set $X \cup Y$ with a fixed order (X, Y) of its parts, i.e., where $X \cap Y = \emptyset$ and $A \subseteq (X \times Y) \cup (Y \times X)$.

Definition 2 Given a dihypergraph H , we construct a directed bigraph B_H as follows. For every vertex v of H , we introduce a vertex in X , which is also called v . For every hyperarc $a = (a^-, a^+)$, we introduce a vertex $a \in Y$. Whenever $v \in a^-$, there is the arc (v, a) in B_H . Whenever $v \in a^+$, there is the arc (a, v) in B_H .

As an example, consider the neighborhood dihypergraph $H = (V, A)$ of the digraph D shown in Figure 2: $V = \{u, v, w, x\}$, $A = \{a_u, a_v, a_w, a_x\}$, where

$$\begin{aligned}
 a_u &= (\{v\}, \emptyset), \\
 a_v &= (\{w\}, \{u, w\}), \\
 a_w &= (\{v, x\}, \{v\}), \text{ and} \\
 a_x &= (\emptyset, \{w\}).
 \end{aligned}$$

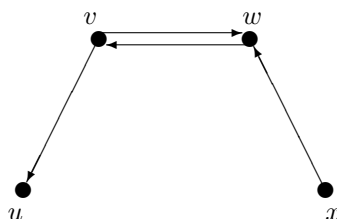


Figure 2: A digraph D .

The directed bigraph B_H of H is shown in Figure 3.

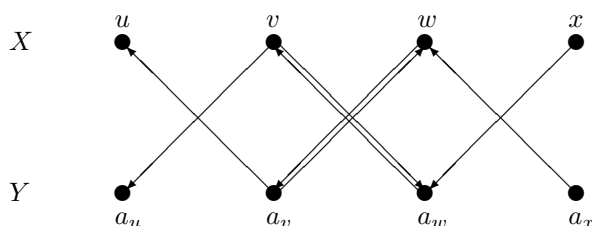


Figure 3: The directed bigraph B_H of H .

Consider a directed bigraph $B = (X, Y, A)$ and an automorphism $\alpha : (X \cup Y) \rightarrow (X \cup Y)$ of the underlying bipartite digraph B , that is for which $(i, j) \in A$ if and only if $(\alpha(i), \alpha(j)) \in A$. The automorphism α is *involutory* if $\alpha(i) = j$ implies $\alpha(j) = i$, that is α^2 is identity, and it is called *switching* if $\alpha(X) = Y$ and $\alpha(Y) = X$. The Digraph Realization Problem for a directed hypergraph H can be equivalently formulated in terms of B_H : Does B_H admit an involutory switching automorphism α such that x and $\alpha(x)$ are non-adjacent for all $x \in X$?

To a directed bigraph $B = (X, Y, A)$ we can associate its X - Y -adjacency matrix $A(B) = (a_{ij}) \in \{0, 1, -1, \pm 1\}^{X \times Y}$ defined by

- $a_{ij} = 0$ if and only if $i \in X, j \in Y, (i, j) \notin A$ and $(j, i) \notin A$,
- $a_{ij} = 1$ if and only if $i \in X, j \in Y, (i, j) \in A$ and $(j, i) \notin A$,
- $a_{ij} = -1$ if and only if $i \in X, j \in Y, (j, i) \in A$ and $(i, j) \notin A$,

- $a_{ij} = \pm 1$ if and only if $i \in X, j \in Y, (i, j) \in A$ and $(j, i) \in A$.

We have

$$A(B_H) = (0 - 10010 \pm 100 \pm 10 - 10010)$$

for the directed bigraph B_H of Figure 3.

3 Orgraph realizations and skew symmetrization

An *orgraph* is an orientation of an undirected graph. In other words, an orgraph is a digraph having no pairs of opposite arcs. Here we consider Decision Problem 1 for orgraphs – the Orgraph Realization Problem.

Theorem 2 *The Orgraph Realization Problem is NP-complete.*

We construct a polynomial-time reduction from the Neighborhood Realization Problem, which is NP-complete by Theorem 1. Let H be an instance to the problem represented as a bigraph $B = (X, Y, E)$. In terms of B , the problem is to recognize whether B has an involutory automorphism α (that is α^2 is identical) which switches the parts ($\alpha(X) = Y$). Without loss of generality, we may assume that all vertex degrees in B are at least three. To satisfy this assumption we can add $i \leq 3$ new vertices into each part, making them adjacent to all vertices in the opposite part.

Now we transform B into a directed bigraph $B' = (X', Y', A)$ by replacing every edge $e = xy \in E$, where $x \in X$ and $y \in Y$, by a directed 6-cycle

$$C^e = (x = x_1^e, y_1^e, x_2^e, y = y_2^e, x_3^e, y_3^e), \tag{1}$$

and put the vertices x_i^e and y_i^e into the parts X' and Y' of B' , respectively, see Figure 4 for an illustration.

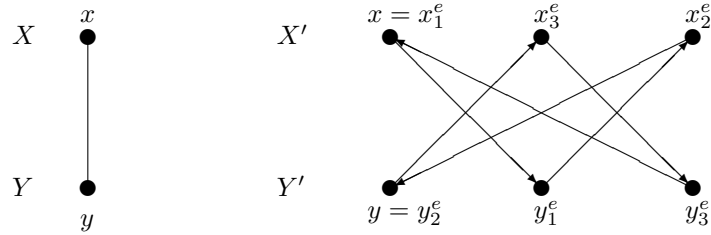


Figure 4: The construction of a directed bigraph $B' = (X', Y', A)$.

The directed bigraph B' represents a dihypergraph H' which is considered as an instance to the Orgraph Realization Problem. In terms of B' , the problem is to recognize whether B' has an involutory automorphism α' which switches the parts X and Y' , and such that x' and $\alpha'(x')$ are always non-adjacent, where $x' \in X'$.

Suppose that B admits an involutory automorphism α that switches the parts X and Y . If some vertices $x \in X$ and $y = \alpha(x) \in Y$ are adjacent, let $a = xy$, then we define $\alpha'(x) = y$, $\alpha'(x_2^a) = y_3^a$ and $\alpha'(x_3^a) = y_1^a$, see the correspondence in Figure 4. Now consider two edges $e = xv$ and $f = uy$ of B such that $y = \alpha(x) \neq v = \alpha(u)$, as it is shown in Figure 5.

The vertices

$$x = x_1^e, y_1^e, x_2^e, v = y_2^e, x_3^e, y_3^e$$

of the directed cycle C^e will be mapped by α' to the vertices

$$y = y_2^f, x_3^f, y_3^f, u = x_1^f, y_1^f, x_2^f$$

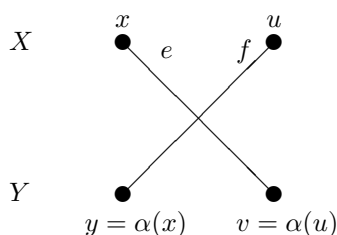


Figure 5: The edges $e = xv$ and $f = uy$ of B .

of the directed cycle C^f , respectively, as it is shown in Figure 6. It is easy to see that α' is an involutory automorphism of B' that switches X' and Y' . Also, x' and $\alpha'(x')$ are non-adjacent for all $x' \in X'$.

Conversely, let α' be an involutory automorphism of B' switching X' and Y' , and such that x' and $\alpha'(x')$ are non-adjacent for all $x' \in X'$. The degree assumption implies that α' pairs the vertices of X with the vertices of Y . Thus, α' induces an involutory bijection α on B that switches X and Y . Finally, α is an automorphism of B . Indeed, let $y = \alpha(x)$ and $v = \alpha(u)$ for some distinct vertices $x, u \in X$. Suppose that $e = xv$ is an edge of B . It is easy to see that the directed 6-cycle C^e can be mapped by α' to another directed 6-cycle as in Figure 6 only. It shows that u and y must be adjacent.

A square matrix $A = (a_{ij})$ is called *skew* if $a_{ij} = -a_{ji}$ for all i and j . In other words, $A = -A^T$, where A^T is the transpose of A . Clearly, all entries on the main diagonal must be zeroes. A square matrix A is *skew-symmetrizable* if it is possible to obtain a skew matrix permuting rows of A .

Decision Problem 2 (Skew-Symmetrization Problem)

Instance: A square $\{0, 1, -1\}$ matrix A .

Question: Is A a skew-symmetrizable matrix?

The Orgraph Realization Problem is essentially the same as the Skew-Symmetrization Problem. Let a dihypergraph H be an instance to the Orgraph Realization Problem. We may assume that $|V(H)| = |A(H)|$. The directed bigraph B of H does not have pairs of opposite arcs (otherwise H has no orgraph realizations). The $\{0, 1, -1\}$ adjacency matrix of B is skew-symmetrizable if and only if $H = \mathcal{N}(D)$ for some orgraph D .

Corollary 2 *The Matrix Skew-Symmetrization Problem is NP-complete.*

It is interesting to study the Matrix Skew-Symmetrization Problem within hereditary classes of orgraphs, in particular for D -free orgraphs.

4 Skew transposability

We write $A \rightarrow B$ if a matrix A can be transformed to a matrix B with row and column permutations. Here we consider the following problem which is related to skew symmetrizability. A square matrix A is *skew-transposable* if $A \rightarrow -A^T$, where A^T is the transpose of A .

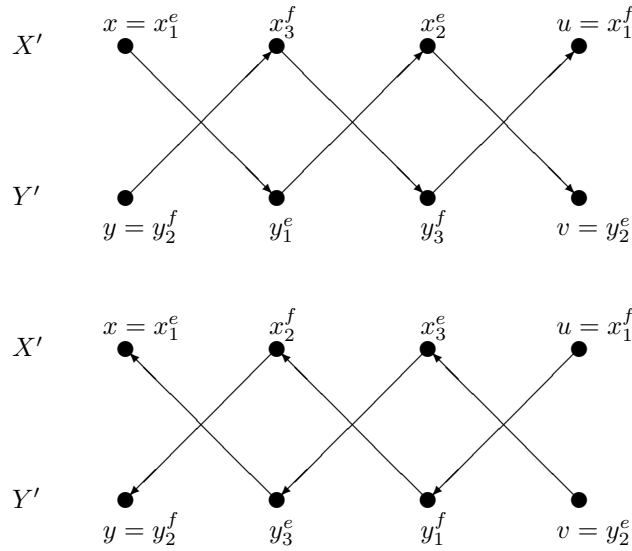


Figure 6: The automorphism α' .

Decision Problem 3 (Skew Transposability Problem)

Instance: A square $\{0, 1, -1\}$ matrix A .

Question: Is A a skew-transposable matrix?

Here is a relation between the two problems.

Proposition 1 Every skew-symmetrizable matrix A is skew-transposable.

By the definition of skew-symmetrizable, there exists a permutation matrix P such that PA is skew-symmetric, that is $PA = -(PA)^T = -A^T P^T$. To show that $A \rightarrow -A^T$, we apply P to the columns of PA : $PAP = -A^T P^T P = -A^T$, meaning that A skew-transposable.

If we represent a square $\{0, 1, -1\}$ matrix A as a directed bigraph $B = (X, Y, A)$, then the matrix $-A^T$ produces the *reversed* directed bigraph $B' = (Y, X, A)$. For example, let

$$A = (10 - 1011 - 100).$$

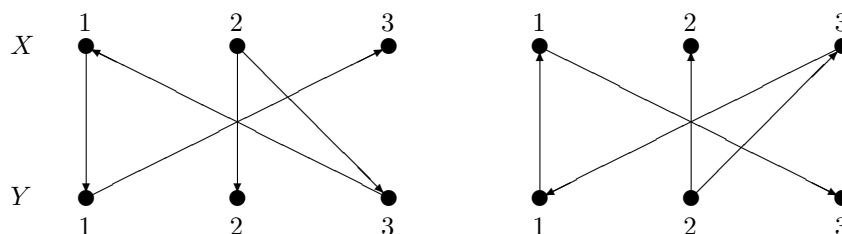
We have

$$-A^T = (-1010 - 101 - 10).$$

The corresponding directed bigraphs B and B' are shown in Figure 7.

Now we clarify the complexity of Decision Problem 3.

Proposition 2 The Skew Transposability Problem is polynomial-time equivalent to graph isomorphism.

Figure 7: The directed bigraphs B and B' .

The Skew Transposability Problem is equivalent to checking whether B and B' are isomorphic, which a particular case of graph isomorphism. Conversely, suppose we want to check isomorphism of graphs G and H . We represent G as a directed bigraph $B_G = (X_G, Y_G, A_G)$, where $X_G = V(G)$, $Y_G = E(G)$, and every edge $e = uv \in E(G)$ produces two arcs (u, e) and (v, e) in B . A similar bigraph $B_H = (X_H, Y_H, A_H)$ is defined for H , and $B'_H = (Y'_H, X'_H, A'_H)$ is obtained by reversing of B_H . Let B is disjoint union of B_G and B'_H . Accordingly, B' is disjoint union of B'_G and B_H . Assuming that both G and H do not have isolated vertices, G and H are isomorphic if and only if B and B' are.

5 Digraph realizations within hereditary classes

Fomin, Kratochvíl, Lokshtanov, Mancini, and Telle [15] studied the Open Neighborhood Realization Problem within hereditary classes.

Definition 3 Let \mathcal{P} be hereditary class of graphs. A \mathcal{P} -realization of a hypergraph H is a graph $G \in \mathcal{P}$ such that $\mathcal{N}(G) = H$. If \mathcal{P} is defined by a unique minimal forbidden induced subgraph H , then a \mathcal{P} -realization is called an H -free realization of H .

Definition 3 is extended to digraphs in a straightforward way.

A *star-like graph* consists of $k \geq 1$ paths $Q_i = (u_0, u_{i1}, u_{i2}, \dots, u_{id_i})$, $i = 1, 2, \dots, k$, having a common vertex u_0 . Here $d_i \geq 0$ for $i = 1, 2, \dots, k$. An example of a star-like graph with $k = 3$, $d_1 = 3$, $d_2 = 4$, and $d_3 = 2$ is shown in Figure 8.

If every connected component of a graph G is star-like, then G is called an *S-graph*. Fomin, Kratochvíl, Lokshtanov, Mancini, and Telle [15] proved the following result in the complementary form (for closed neighborhood hypergraphs).

Theorem 3 If H is not an *S-graph*, then it is NP-hard to decide whether a given hypergraph has an H -free realization.

Theorem 3 can be easily extended to \mathcal{P} -realizations, where \mathcal{P} is a hereditary class with a finite set $Z(\mathcal{P})$ of minimal forbidden induced subgraphs.

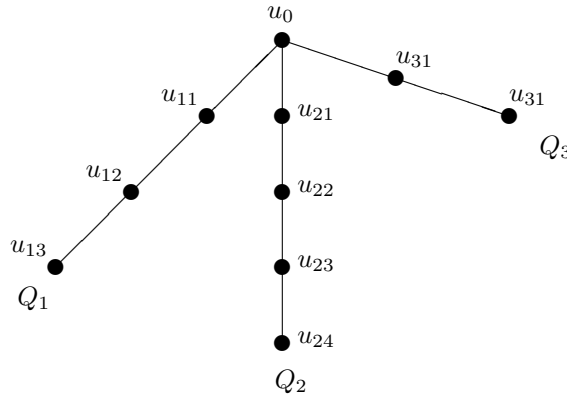


Figure 8: An example of a star-like graph.

Theorem 4 *If $Z(\mathcal{P})$ is a finite set and it does not contain an S -graph, then it is NP-hard to decide whether a given hypergraph has a \mathcal{P} -realization.*

If H is an S -graph, then complexity of the H -free realization problem is unknown, except the following polynomial-time solvable cases: $H \in \{\overline{P}_1, \overline{P}_2, \overline{P}_3, \overline{P}_4, \overline{C}_3, \overline{C}_4\}$, where P_k and C_k are the path and the cycle with k vertices, and \overline{G} is the complement of G , see Fomin, Kratochvíl, Lokshtanov, Mancini, and Telle [15].

We are going to extend Theorem 3 and Theorem 4 to digraphs.

A *star-like digraph of type 1* is obtained from a star-like graph G if we replace every edge $uv \in E(G)$ by the corresponding pair $(u, v), (v, u)$ of opposite arcs. A *star-like digraph of type 2* consists of $k \geq 1$ directed paths

$$Q_i = (u_0, u_{i1}, u_{i2}, \dots, u_{id_i}),$$

$i = 1, 2, \dots, k$, having a common vertex u_0 , and of $l \geq 0$ directed paths

$$R_j = (v_{j1}, v_{j2}, \dots, v_{je_j}, u_0),$$

$j = 1, 2, \dots, l$, having a common vertex u_0 . Here $d_i \geq 0$ and $e_j \geq 0$ for all i and j . An example of a star-like graph with $k = 3, d_1 = 3, d_2 = 4, d_3 = 2, l = 2, e_1 = 3$ and $e_2 = 2$ is shown in Figure 9.

If every weakly connected component of a digraph D is a star-like digraph of type i , then D is called an S_i -digraph, $i = 1, 2$.

Theorem 5 *If a digraph D has at least one arc, then it is NP-hard to decide whether a given dihypergraph H has a D -free realization.*

First we apply Theorem 3 to a *symmetric* dihypergraph H , that is $a^- = a^+$ for every hyperarc (a^-, a^+) of H .

Property 1 *If D is not an S_1 -digraph, then it is NP-hard to decide whether a symmetric dihypergraph H has a D -free realization.*

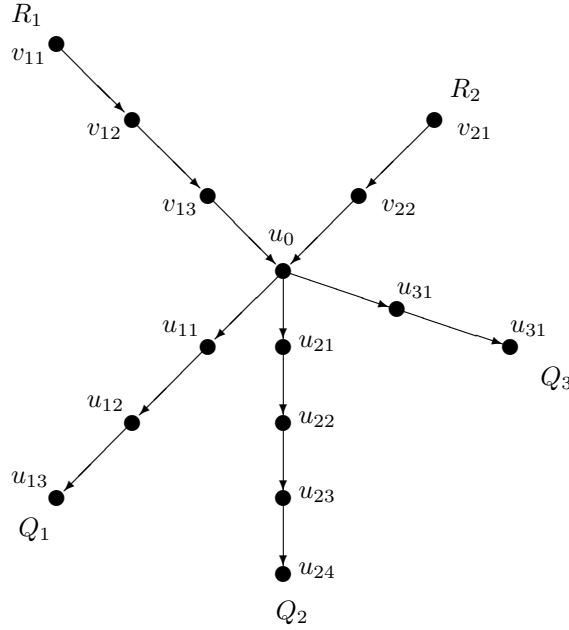


Figure 9: An example of a star-like digraph of type 2.

A digraph is *symmetric* if (u, v) is an arc if and only if (v, u) is an arc. Essentially, a symmetric digraph is an undirected graph. Clearly, every realization of a symmetric dihypergraph is a symmetric digraph, and Theorem 3 implies the result, since D is not an S_1 -digraph.

Now we consider S_2 -digraphs.

Property 2 *If D is not an S_2 -digraph, then it is NP-hard to decide whether a given directed hypergraph has a D -free realization.*

We modify the proof of Theorem 2 in the following way. Instead of a directed 6-cycle C^e for an edge $e = xy$ as in (1), we introduce a $(4t + 2)$ -cycle C^e

$$C^e = (x = x_1^e, y_1^e, x_2^e, y_2^e, \dots, x_t^e, y = y_t^e, \dots, x_{2t+1}^e, y_{2t+1}^e) \quad (2)$$

for a fixed $t \geq 1$. The resulting dihypergraph and directed bigraph are denoted by H' and B' , respectively. We shall specify t so that every realization of H' does not contain the forbidden induced subdigraph D . Let t_1 be the minimum length of a cycle (not necessarily directed) in D . If D is acyclic then $t_1 = \infty$. A *knot vertex* of D is a vertex u such that either

- $|N^-(u)| + |N^+(u)| \geq 3$, or
- $|N^-(u)| = 2$, or
- $|N^+(u)| = 2$.

Let t_2 be the minimum length of a path (not necessarily directed) in D that connects two knot vertices in D . If D does not have such paths, then $t_2 = \infty$. At least one of t_1 and t_2 is finite, since D is not an S_2 -digraph. It is sufficient to take $t = \min\{t_1, t_2\}$.

Property 1 and Property 2 show that the problem is NP-hard unless D is both an S_1 -digraph and an S_2 -digraph. But it is possible only if D does not have arcs.

Let O_n be an arcless digraph of order n .

Open Problem 1 *How hard is to decide whether a given directed hypergraph has an O_n -free realization, $n \geq 3$?*

For $n \leq 2$, the problem is trivially polynomial-time solvable.

6 Matrix complementation

Here we consider another interesting problem related to 0 – 1 matrices. Let $A = (a_{ij})$ be an $m \times n$ matrix with $a_{ij} \in \{0, 1\}$, $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$. The *complement* of A is the matrix $\bar{A} = (\bar{a}_{ij})$ defined by: $\bar{a}_{ij} = 1 - a_{ij}$ for all i and j . Recall that $A \rightarrow B$ means that a matrix A can be transformed to a matrix B with row and column permutations.

Decision Problem 4 (Matrix Complementation Problem)

Instance: A 0 – 1 matrix A .

Question: *Does $A \rightarrow \bar{A}$ hold?*

As an example, consider the matrix

$$A = (101001).$$

Permuting row 1 and row 2, we obtain

$$(001101).$$

Now, permutation of column 2 and column 3 gives

$$(010110) = \bar{A},$$

therefore $A \rightarrow \bar{A}$.

We show that the Matrix Complementation Problem is polynomial-time equivalent to graph isomorphism. One can mention a related result of McCarthy and McKay [20] which says that the problem $A \rightarrow A^T$, where A is a square 0 – 1 matrix A and A^T is the transpose of A , is also polynomial-time equivalent to graph isomorphism.

An obvious necessary condition for $A \rightarrow \bar{A}$ is that $A_0 = A_1$, where A_k denotes the total number of entries $a_{ij} = k$ in A . However, this condition is not sufficient. For example, it is impossible to get \bar{A} from the matrix

$$A = (100100110011),$$

where $A_0 = A_1 = 6$. Indeed, permuting columns of A , one can obtain the following six matrices:

$$(100100110011), (100100101011), (010010110101), (010010011101), (001001101110), (001001011110),$$

and, unlike \bar{A} , no one of them has two rows (011). Thus, $A \rightarrow \bar{A}$ does not hold.

Theorem 6 *The Matrix Complementation Problem and the Graph Isomorphism Problem are polynomial-time equivalent.*

First we represent A and \bar{A} as bigraphs $B = (X, Y, E)$ and $B' = (X', Y', E')$, respectively. The bigraphs B and B' are *isomorphic* if there are bijections $\alpha : X \leftrightarrow X'$ and $\beta : Y \leftrightarrow Y'$ such that $(i, j) \in E$ if and only if $(\alpha(i), \beta(j)) \in E'$. The corresponding recognition problem is called *Bigraph Isomorphism*.

Fact 1 $A \rightarrow \bar{A}$ holds if and only if the bigraphs B and B' are isomorphic.

Indeed, a permutation α of rows and a permutation β of columns is nothing but an isomorphism of corresponding bigraphs.

The *bi-complement* of B is the bigraph $\bar{B} = (X, Y, \bar{E})$, where

$$\bar{E} = \{xy : x \in X, y \in Y, xy \notin E\}.$$

Clearly, B' is isomorphic to \bar{B} . A bigraph is *self-bi-complementary* if B and \bar{B} are isomorphic, see Bhawe and Raghunathan [6]. In this terminology, Fact 1 says that $A \rightarrow \bar{A}$ holds if and only if B is a self-bi-complementary bigraph. Recognition of self-bi-complementary bigraphs is a particular case of the Bigraph Isomorphism Problem, therefore the Matrix Complementation Problem is not harder than graph isomorphism.

Fact 2 The Graph Isomorphism Problem is polynomial-time reducible to recognition of self-bi-complementary bigraphs.

Let G and H be an instance to the Graph Isomorphism Problem. Without loss of generality, we may assume that $|V(G)| = |V(H)| = n$, $|E(G)| = |E(H)| = m$ (otherwise G and H are not isomorphic) and both G and H do not have isolated vertices (otherwise we add a dominating vertex to each of them obtaining an equivalent instance).

We subdivide every edge of G and H with a new vertex, and denote the resulting graphs by G' and H' , respectively. G' can be considered as a bigraph having $V(G)$ as its X -part (old vertices) and the set of $|E(G)|$ new vertices as its Y -part. Similar situation takes place for H' . Now we use the graphs G' and H' to construct a bigraph $B = (X, Y, E)$ such that $G \cong H$ if and only if B is self-bi-complementary. For that, we take disjoint copies of G' and \bar{H}' [the bi-complement of H'], and introduce all edges between the X -part of G' and and the Y -part of \bar{H}' . Figure 10 illustrates the construction.

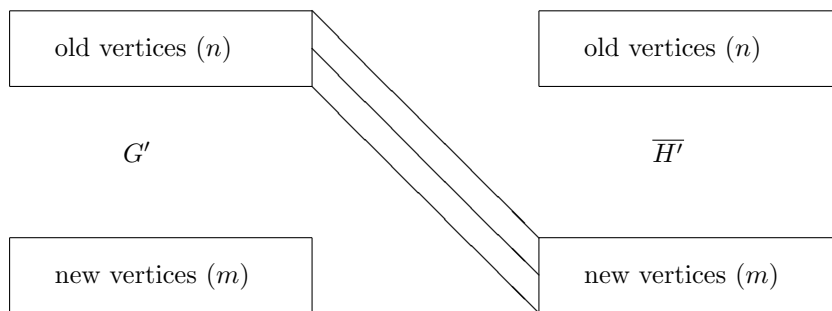


Figure 10: The construction of B .

The bi-complement \bar{B} of B is shown in Figure 11, where \bar{G}' and H' are the bi-complements of G' and H' , respectively, and all edges between the X -part H' of and the Y -part of \bar{G}' are included.

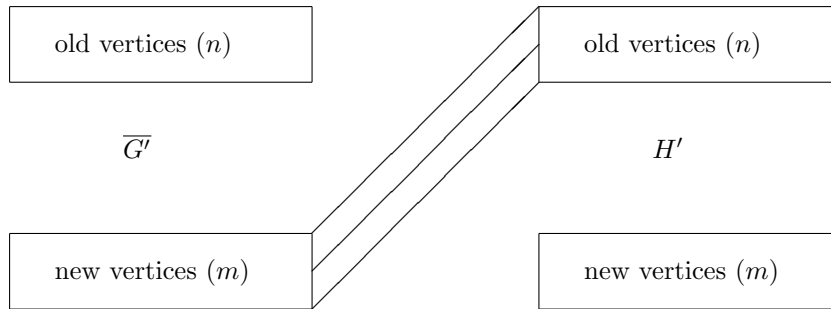


Figure 11: The bi-complement \overline{B} of B .

If we have an isomorphism $\phi : V(G) \rightarrow V(H)$, then we can obviously extend ϕ to isomorphisms of G' and H' , and $\overline{H'}$ and $\overline{G'}$. In turn, they induce an isomorphism of the bigraphs B and \overline{B} .

Conversely, let α, β be an isomorphism of B and \overline{B} . The assumptions imply that $\deg_B u \geq m+1 > \deg_{\overline{B}} v$ for all old vertices u, v of G' . It shows that α transforms the old vertices of G' to the old vertices of H' . Similarly, $\deg_B u = 2 < n+2 \leq \deg_{\overline{B}} v$ for all new vertices u, v of G' . Hence β transforms the new vertices of G' to the new vertices of H' . As a result, we obtain an isomorphism of G' and H' which induces an isomorphism of G and H .

Now the result follows from Fact 1 and Fact 2.

Fact 2 is similar to a known result of Colbourn and Colbourn [14, 12] that recognizing whether a graph is self-complementary is polynomially equivalent to the graph isomorphism problem. The Matrix Complementation Problem can be viewed as a particular case of the following *Matrix Negation Problem* (if we replace 0 by -1): Given a matrix A over a set of integers, whether $A \rightarrow -A$. It is not hard to show that the Matrix Negation Problem is polynomial-time equivalent to graph isomorphism.

7 Tournament realizations and anti-symmetrization

A *tournament* is an orientation of a complete undirected graph. Decision Problem 1 for tournaments is trivial. However, Aigner and Triesch [2] proposed an interesting variant of the problem. Given a digraph $D = (V, A)$, define the *(+)-neighborhood hypergraph*, $H = \mathcal{N}^+(D)$, by $V(H) = V$ and $E(H) = \{N^+(u) : u \in V\}$.

Decision Problem 5 (Digraph (+)-Realization Problem)

Instance: A hypergraph H .

Question: Does $H = \mathcal{N}^+(D)$ hold for some digraph D ?

This problem is simple in general: Aigner and Triesch [2] noted that it is equivalent to finding a perfect matching in a bipartite graph. But they were unable to solve Decision Problem 5 for tournaments.

We represent a hypergraph H as an (undirected) bigraph $B = (X, Y, E)$. The problem is to find an involutory switching automorphism α such that x and $\alpha(x)$ are always non-adjacent, and $x \in X$ is adjacent to $\alpha(x') \in Y$ if and only if the vertices $x' \in X$ and $\alpha(x) \in Y$ are non-adjacent. Illustrations for the oriented triple and the transitive triple are given in Figure 12 and Figure 13, respectively.

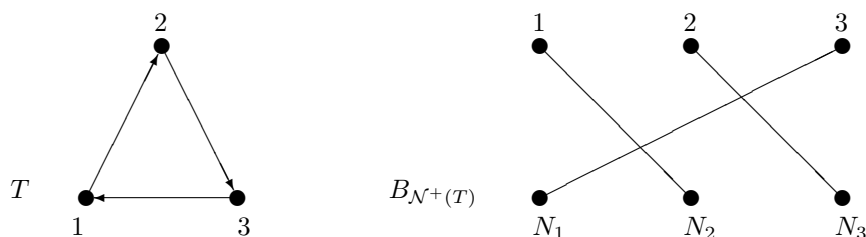


Figure 12: An illustration for the oriented triple.

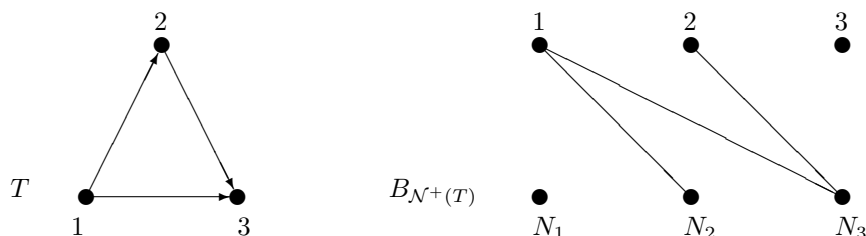


Figure 13: An illustration for the transitive triple.

To a bigraph $B = (X, Y, E)$ we can associate its X - Y -adjacency matrix $A(B) = (a_{ij}) \in \{0, 1\}^{X \times Y}$ defined by $a_{ij} = 1$ if and only if $(i, j) \in E$. Conversely, any 0–1 matrix $A = (a_{ij})$ can be viewed as the X - Y adjacency matrix $A = A(B)$ of a corresponding bigraph $B = (X, Y, E)$, where X is the set of row indices of A , Y is the set of column indices of A , and $(i, j) \in E$ if and only if $a_{ij} = 1$. Here are the adjacency matrices of the bigraphs of Figure 12 and Figure 13, respectively:

$$(010001100) \quad \text{and} \quad (011001000).$$

Now we reformulate the problem in terms of square 0–1 matrices as follows. Does a given 0–1 square matrix A admits a permutation of rows such that the resulting matrix B has the properties:

(all-0 diagonal) $b_{ii} = 0$ for all i , and

(anti-symmetry) $b_{ij} \neq b_{ji}$ for all $i \neq j$?

It is called the *Matrix Anti-Symmetrization Problem*.

Conjecture 1 *The Matrix Anti-Symmetrization Problem is NP-hard.*

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