

A REVIEW OF THE S-PROCEDURE

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1 Overview.

The s-procedure is an instrumental tool in control theory and robust optimization analysis. It is also used in semi-definite relaxation of quadratic programming. In this survey, we first define the s-procedure and prove the losslessness theorem. And in the second part, we introduce the SDP-relaxation of quadratic programming, which is an application of the s-procedure.

2 S-lemma.

The s-procedure is a tool verifying the non-negativity of a quadratic function under quadratic constraints. We begin with a first result dealing the quadratic optimization problem under a single quadratic inequality constraints, proved by Yakubovich (1971):

Theorem 1 (S-lemma) $A, B \in \mathbb{S}_n$, we assume $x^T A x > 0$ is feasible, then we have

$$(x^T A x \geq 0 \Rightarrow x^T B x \geq 0)$$

if and only if

$$\exists \lambda \geq 0, B \succeq \lambda A$$

Proof: Sufficiency is clear. Assume $(x^T A x \geq 0 \Rightarrow x^T B x \geq 0)$. Let

$$S = \{(x^T A x, x^T B x) \mid x \in \mathbb{R}^n\},$$

and

$$\mathbf{U} = \{(\mathbf{u}_1, \mathbf{u}_2) | \mathbf{u}_1 \geq 0, \mathbf{u}_2 < 0\}.$$

S is convex by Dines' theorem (1941), and \mathbf{U} is a convex cone, as $S \cap \mathbf{U} = \emptyset$, there exists a separating hyperplane $\mathbf{c} = (\mathbf{c}_1, \mathbf{c}_2)$ such that $(\mathbf{c}, \mathbf{s}) \leq 0, \forall \mathbf{s} \in S$ and $(\mathbf{c}, \mathbf{u}) \geq 0, \forall \mathbf{u} \in \mathbf{U}$, because

Lemma 2 *Two nonempty convex subset of \mathbb{R}^n can be properly separated by a hyperplane if and only if their relative interiors are disjoint.*

For any $\alpha > 0, (1, -\alpha) \in \mathbf{U}$, ie, $\mathbf{c}_1 \geq \alpha \mathbf{c}_2$, so $\mathbf{c}_1 \geq 0$. $\bar{\mathbf{x}}$ is such that $\bar{\mathbf{x}}^T \mathbf{A} \bar{\mathbf{x}} > 0$, then

$$\mathbf{c}_1 \bar{\mathbf{x}}^T \mathbf{A} \bar{\mathbf{x}} + \mathbf{c}_2 \bar{\mathbf{x}}^T \mathbf{B} \bar{\mathbf{x}} \leq 0$$

implies that $\mathbf{c}_2 < 0$, as $\mathbf{c}_1 \mathbf{c}_2 \neq 0$. Therefore,

$$\mathbf{x}^T \mathbf{B} \mathbf{x} \geq -\frac{\mathbf{c}_1}{\mathbf{c}_2} \mathbf{x}^T \mathbf{A} \mathbf{x}, \forall \mathbf{x}.$$

■

3 S-procedure.

The S-procedure is the generation of S-lemma. Let

$$\begin{aligned} \sigma_k : V &\rightarrow \mathbb{R} \\ \sigma_k(\mathbf{y}) &= \mathbf{y}^T \mathbf{Q}_k \mathbf{y} + 2\mathbf{s}_k^T \mathbf{y} + r_k, \end{aligned}$$

for $k = 0, 1, \dots, N$, be quadratic fuctions defined on a linear vector space V and consider the two conditions:

- S_1 : For all $\mathbf{y} \in V$,

$$(\sigma_k(\mathbf{y}) \geq 0, \forall k = 1, \dots, N) \Rightarrow \sigma_0(\mathbf{y}) \geq 0.$$

- S_2 : There exists $\tau_k \geq 0, k = 1, \dots, N$ such that

$$\sigma_0(\mathbf{y}) - \sum_{k=1}^N \tau_k \sigma_k(\mathbf{y}) \geq 0, \forall \mathbf{y} \in V.$$

The S-procedure is the method verifying S_1 using S_2 . There are some reasons for which it is generally hard to verify S_1 :

1. σ_0 is not a convex function in general.
2. The constraint set

$$\{\mathbf{y} | \sigma_k(\mathbf{y}) \geq 0, \forall k = 1, \dots, N\}$$

is not convex in general.

However S_2 is easy to verify. In fact,

$$\begin{aligned}
 S_2 &\Leftrightarrow \exists \tau_k \geq 0 \text{ s.t. } \sigma_0(\mathbf{y}) - \sum_{k=1}^N \tau_k \sigma_k(\mathbf{y}) \geq 0, \forall \mathbf{y} \\
 &\Leftrightarrow \exists \tau_k \geq 0 \text{ s.t. } \begin{pmatrix} \mathbf{y} \\ 1 \end{pmatrix}^T \begin{pmatrix} Q_0 + \sum_{k=1}^N \tau_k Q_k & s_0 + \sum_{k=1}^N \tau_k s_k \\ s_0^T + \sum_{k=1}^N \tau_k s_k^T & r_0 + \sum_{k=1}^N \tau_k r_k \end{pmatrix} \begin{pmatrix} \mathbf{y} \\ 1 \end{pmatrix} \geq 0, \forall \mathbf{y} \\
 &\Leftrightarrow \exists \tau_k \geq 0 \text{ s.t. } z^T \begin{pmatrix} Q_0 + \sum_{k=1}^N \tau_k Q_k & s_0 + \sum_{k=1}^N \tau_k s_k \\ s_0^T + \sum_{k=1}^N \tau_k s_k^T & r_0 + \sum_{k=1}^N \tau_k r_k \end{pmatrix} z \geq 0, \forall z = \begin{pmatrix} \mathbf{y} \\ \delta \end{pmatrix} \text{ s.t. } \delta \neq 0 \\
 &\Leftrightarrow \exists \tau_k \geq 0 \text{ s.t. } z^T \begin{pmatrix} Q_0 + \sum_{k=1}^N \tau_k Q_k & s_0 + \sum_{k=1}^N \tau_k s_k \\ s_0^T + \sum_{k=1}^N \tau_k s_k^T & r_0 + \sum_{k=1}^N \tau_k r_k \end{pmatrix} z \geq 0, \forall z \\
 &\Leftrightarrow \exists \tau_k \geq 0 \text{ s.t. } \begin{pmatrix} Q_0 & s_0 \\ s_0^T & r_0 \end{pmatrix} + \sum_{k=1}^N \tau_k \begin{pmatrix} Q_k & s_k \\ s_k^T & r_k \end{pmatrix} \succeq 0.
 \end{aligned}$$

Obviously, $S_2 \Rightarrow S_1$. In general $S_1 \not\Rightarrow S_2$, however there are some special cases when $S_1 \Leftrightarrow S_2$. As in Theorem 1, we need some regularity on functions σ_k :

Definition 3 *The constraint $\sigma_k(\mathbf{y}) \geq 0, \forall k = 1, \dots, N$ is said to be regular if there exists $\mathbf{y}^* \in \mathcal{V}$ such that $\sigma_k(\mathbf{y}^*) > 0, \forall k = 1, \dots, N$.*

Then we have the following theorem:

Theorem 4 (Yakubovich) *Let*

$$\mathcal{K} = \{(\sigma_0(\mathbf{y}), \sigma_1(\mathbf{y}), \dots, \sigma_N(\mathbf{y})) | \mathbf{y} \in \mathcal{V}\},$$

$$\mathcal{N} = \{(n_0, n_1, \dots, n_N) | n_0 < 0, n_k > 0, k = 1, \dots, N\}.$$

If $(\mathcal{K} \cap \mathcal{N} = \emptyset) \Rightarrow (\text{conv}(\mathcal{K}) \cap \mathcal{N} = \emptyset)$, then under the regularity on constraint, the S -procedure is lossless, i.e., $S_1 \Leftrightarrow S_2$.

Proof: Analogy of the proof of Theorem 1. ■

Remark 5 *If \mathcal{K} is convex, the S -procedure is lossless.*

Corollary 6 (S-procedure Losslessness in linear case) *In the linear case, where $\sigma_k(\mathbf{y}) = s_k^T \mathbf{y} + r_k$, under the regularity of constraint, the S -procedure is lossless.*

Corollary 7 (S-procedure Losslessness in the case of one quadratic constraint)

1. *Assume σ_1 is regular, then the following conditions*

- $S_1: \forall \mathbf{y} \in \mathcal{V}, (\sigma_1(\mathbf{y}) \geq 0) \Rightarrow (\sigma_0(\mathbf{y}) \geq 0)$

- S_2 : There exists $\tau \geq 0$ such that

$$\begin{pmatrix} Q_0 & s_0 \\ s_0^\top & r_0 \end{pmatrix} - \tau \begin{pmatrix} Q_1 & s_1 \\ s_1^\top & r_1 \end{pmatrix} \succeq 0$$

are equivalent.

2. Assume $\sigma_1(\mathbf{y}) = \mathbf{y}^\top Q_1 \mathbf{y}$ is regular, then the following conditions

- S_1 : $\forall \mathbf{y} \in \mathcal{V}, (\mathbf{y}^\top Q_1 \mathbf{y} \geq 0, \mathbf{y} \neq 0) \Rightarrow (\mathbf{y}^\top Q_0 \mathbf{y} > 0)$
- S_2 : There exists $\tau \geq 0$ such that $Q_0 - \tau Q_1 \succ 0$

are equivalent.

Proof:

1. \mathcal{K} is convex.
2. Let

$$\epsilon = \min_{\{\mathbf{y} | \mathbf{y}^\top \mathbf{y} = 1, \mathbf{y}^\top Q_1 \mathbf{y} \geq 0\}} \mathbf{y}^\top Q_0 \mathbf{y}$$

then $\mathbf{y}^\top (Q_0 - \epsilon I_{\dim(\mathcal{V})}) \mathbf{y} \geq 0$ for all \mathbf{y} such that $\mathbf{y}^\top Q_1 \mathbf{y} \geq 0$. Thus $Q_0 - \epsilon I_{\dim(\mathcal{V})} - \tau Q_1 \succeq 0$.

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4 SDP Relaxation of Quadratically Constrained Quadratic Programming.

We apply the S-procedure to a general case of quadratic programming problems

$$\begin{aligned} \min \quad & \mathbf{y}^\top Q_0 \mathbf{y} + 2s_0^\top \mathbf{y} + r_0 \\ \text{s.t.} \quad & \mathbf{y}^\top Q_k \mathbf{y} + 2s_k^\top \mathbf{y} + r_k \geq 0, \quad k = 1, \dots, N. \end{aligned} \tag{1}$$

This class of optimization problem includes also integer valued constraints, since $x_i^2 - x_i = 0$ is a valid constraint. Thus the problem is generally NP-hard unless the quadratic functions are convex. Lower bound on the optimization problem (1) can be obtained using the following semidefinite relaxation

$$\begin{aligned} \min \quad & \mathbf{Tr}(Q_0 Y) + 2s_0^\top \mathbf{y} + r_0 \\ \text{s.t.} \quad & \mathbf{Tr}(Q_k Y) + 2s_k^\top \mathbf{y} + r_k \geq 0, \quad k = 1, \dots, N, \\ & \begin{pmatrix} Y & \mathbf{y} \\ \mathbf{y}^\top & 1 \end{pmatrix} \succeq 0. \end{aligned} \tag{2}$$

To see that (2) is a lower bound of (1) we notice that $\begin{pmatrix} Y & \mathbf{y} \\ \mathbf{y}^\top & 1 \end{pmatrix} \succeq 0$ is equivalent to $Y \succeq \mathbf{y} \mathbf{y}^\top$ by Schur complement formula. Since $\mathbf{Tr}(Q_k \mathbf{y} \mathbf{y}^\top) = \mathbf{y}^\top Q_k \mathbf{y}$ we see that (2) must have a lower optimal value than (1). And for the case of single constraint, we have

Proposition 8 *Assume that $\mathbf{y}^\top \mathbf{Q}_1 \mathbf{y} + 2\mathbf{s}_1^\top \mathbf{y} + r_1 \geq 0$ is a regular constraint and that the two optimization problems are strictly feasible, then the following optimization problems have the same objective values*

1. Quadratic programming

$$\begin{aligned} \min \quad & \mathbf{y}^\top \mathbf{Q}_0 \mathbf{y} + 2\mathbf{s}_0^\top \mathbf{y} + r_0 \\ \text{s.t.} \quad & \mathbf{y}^\top \mathbf{Q}_1 \mathbf{y} + 2\mathbf{s}_1^\top \mathbf{y} + r_1 \geq 0. \end{aligned} \tag{3}$$

2. Eigen value for LMI

$$\begin{aligned} \max \quad & \gamma \\ \text{s.t.} \quad & \begin{pmatrix} \mathbf{Q}_0 - \lambda \mathbf{Q}_1 & \mathbf{s}_0 - \lambda \mathbf{s}_1 \\ \mathbf{s}_0^\top - \lambda \mathbf{s}_1^\top & r_0 - \lambda r_1 - \gamma \end{pmatrix} \succeq 0 \end{aligned} \tag{4}$$

3. Semidefinite relaxation of (3) and dual of (4)

$$\begin{aligned} \min \quad & \mathbf{Tr}(\mathbf{Q}_0 \mathbf{Y}) + 2\mathbf{s}_0^\top \mathbf{y} + r_0 \\ \text{s.t.} \quad & \mathbf{Tr}(\mathbf{Q}_1 \mathbf{Y}) + 2\mathbf{s}_1^\top \mathbf{y} + r_1 \geq 0 \\ & \begin{pmatrix} \mathbf{Y} & \mathbf{y} \\ \mathbf{y}^\top & 1 \end{pmatrix} \succeq 0. \end{aligned} \tag{5}$$

Proof: The direct consequence of Corollary 7. ■

References

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