

DIMACS/CCICADA Workshop on Stochastic Networks: Reliability, Resiliency, and Optimization

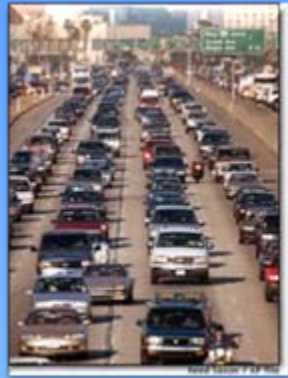
Heterogeneous models for nonlinear flows on networks

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Vehicular Traffic

Lighthill-Witham-Richards model:

$$\rho_t + (\rho v(\rho))_x = 0, \quad \rho \in [0, \rho_{max}] \text{ density of cars}$$


Irrigation Channels

De Saint-Venant equation:

$$\begin{cases} H_t + (Hv)_x = 0 \\ v_t + [\frac{1}{2}v^2 + gH]_x = 0 \end{cases}$$

H water height, v velocity, g gravity



Gas pipelines

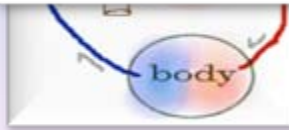
Isothermal Euler with friction:

$$\begin{cases} \rho_t + (\rho u)_x = 0 \\ (\rho u)_t + (\rho u^2 + a^2 \rho)_x = -f_g \frac{\rho u |\rho u|}{2D\rho} \end{cases}$$

ρ density, u velocity

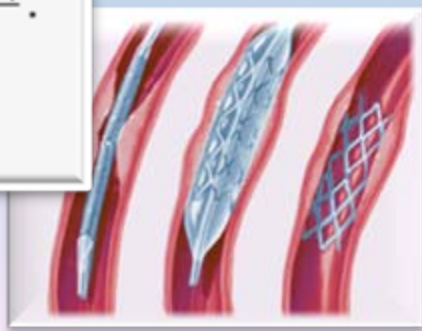
one processing rate

Real fluids



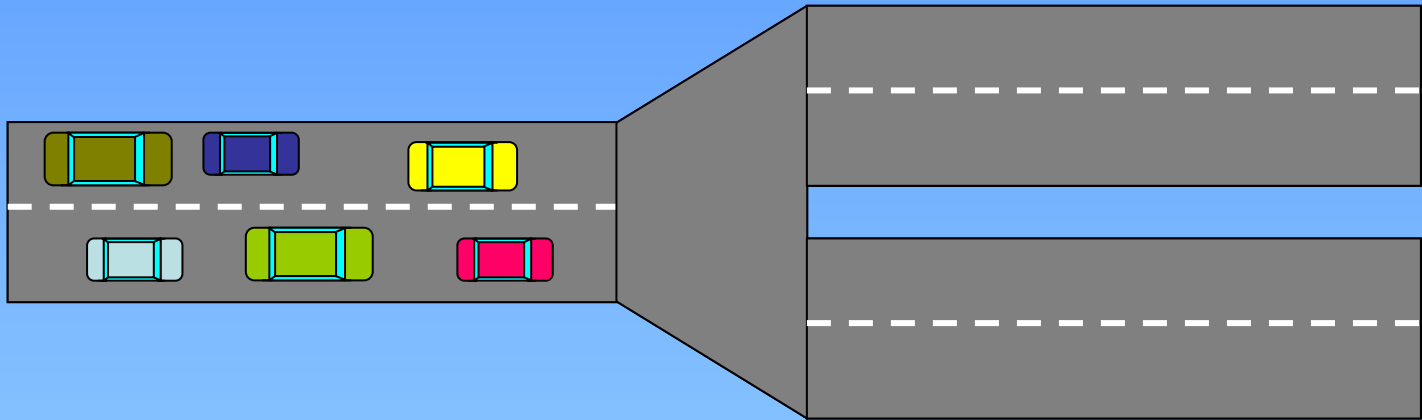
Blood circulation

Bio-Medical



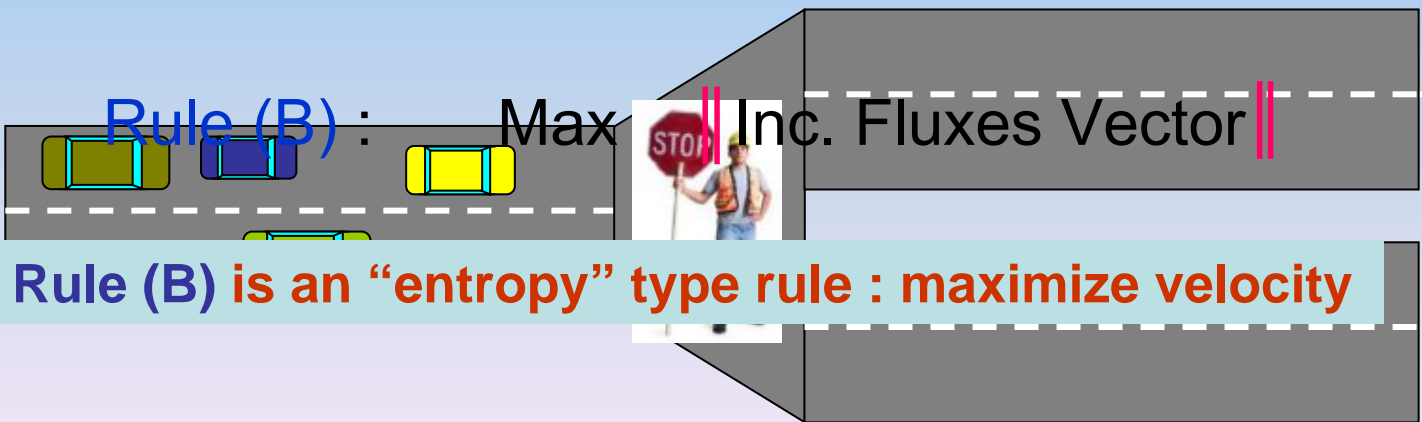
Vascular stents

Dynamics at junctions



Rule (A) : Out. Fluxes Vector = $A \cdot$ Inc. Fluxes Vector

Traffic distribution matrix $A = (\alpha_{ji})$, $0 < \alpha_{ji} < 1$, $\sum_j \alpha_{ji} = 1$



Rule (B) : Max || Inc. Fluxes Vector ||

Rule (B) is an "entropy" type rule : maximize velocity

Integration of models and scales

Car trajectories and moving bottlenecks
Goettlich-Herty-Klar supply chain model

Mixed ODE-PDE model

$$\begin{cases} \partial_t \rho + \dots \\ \rho_t + f(x, y(t), \rho)_x = 0 \end{cases}$$

$$\partial_t \rho_j(x, t) + \partial_x \min\{\mu_j, \frac{L_j}{T_j} \rho_j(x, t)\} = 0 \quad \forall x \in [a_j, b_j], t \in \mathbb{R}^+$$

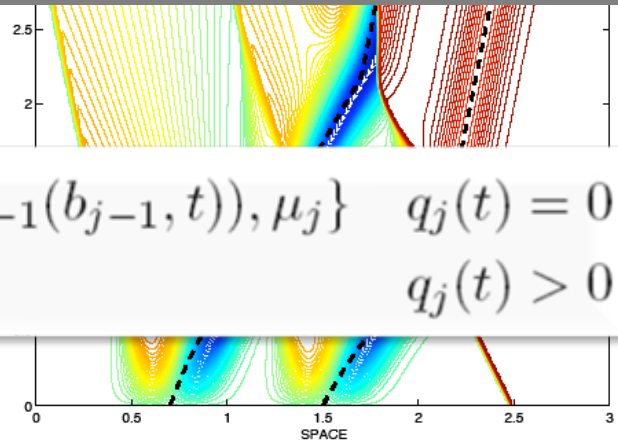
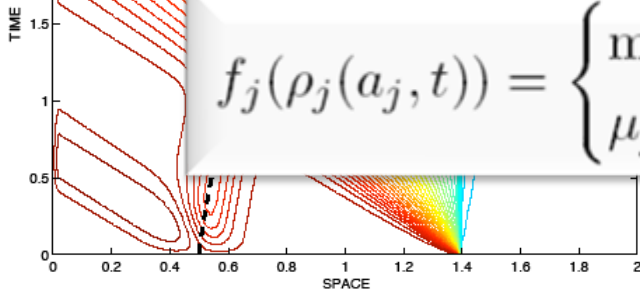
Queue

Processor j

Queue buffer occupancy change is given by the difference between incoming and outgoing flux

$$\partial_t q_j(t) = f_{j-1}(\rho_{j-1}(b_{j-1}, t)) - f_j(\rho_j(a_j, t))$$

$$f_j(\rho_j(a_j, t)) = \begin{cases} \min\{f_{j-1}(\rho_{j-1}(b_{j-1}, t)), \mu_j\} & q_j(t) = 0 \\ \mu_j & q_j(t) > 0 \end{cases}$$

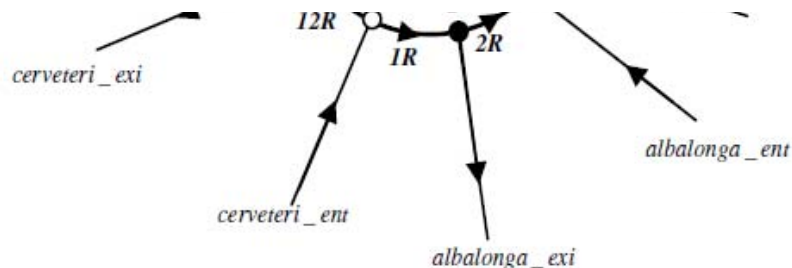
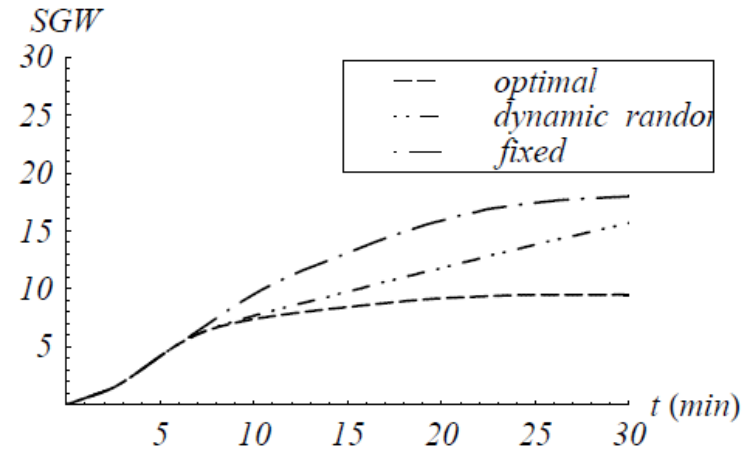
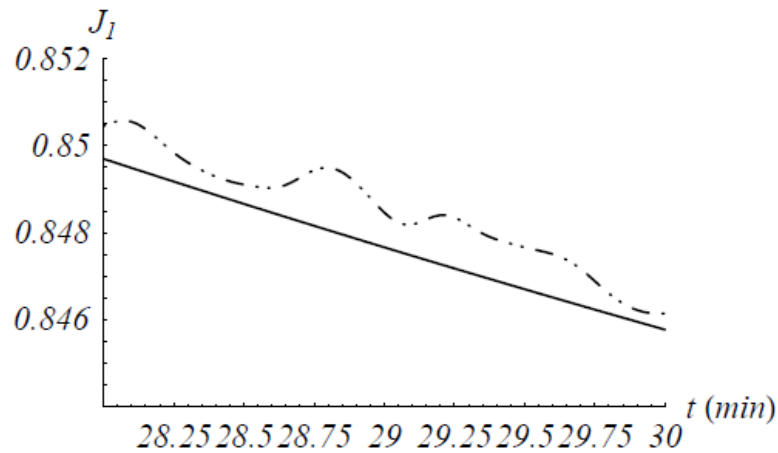


Optimization of vehicular traffic

$$J_1(t) = \sum_i \int_{I_i} v(\rho_i(t, x)) dx,$$

$$J_2(t) = \sum_i \int_{I_i} \frac{1}{v(\rho_i(t, x))} dx.$$

$$SGW = \int_0^T \int_{\cup I_i} |Dv(\rho)| dt dx.$$



Via Parmenide ∪ Via Parmenide

Optimal control for supply chains

$$\begin{cases} \partial_t \rho_j(x, t) + \partial_x \min \{ \mu_j, v_j \rho_j(x, t) \} = 0 & j = 1, \dots, N, \\ \dot{q}_j(t) = f_{j-1}(\rho_{j-1}(b_{j-1}, t)) - f_j^{inc} & j = 2, \dots, N, \\ \rho_1(a_1, t) = u(t) \\ \rho_j(x, 0) = \rho_{j,0}(x) & j = 1, \dots, N, \\ q_j(x, 0) = q_{j,0} & j = 2, \dots, N, \end{cases}$$

$$J(u) = \sum_{j=1}^n \int_0^T q_j(t) dt + \int_0^T [v_N \cdot \rho_N(b_N, t) - \psi(t)]^2 dt \doteq J_1(u) + J_2(u),$$

Existence of solutions

Take minimizing sequence: compactness by Helly and Ascoli Arzela' Theorem.

$$q_n \rightarrow q \text{ in } C^0, \text{ thus } J_1(u_n) \rightarrow J_1(u)$$

$$\begin{aligned} & \int_0^T \left((v_N \cdot \rho_N^n(b_N, t) - \psi(t))^2 - (v_N \cdot \rho_N(b_N, t) - \psi(t))^2 \right) dt = \\ & \int_0^T \left((v_N)^2 \left((\rho_N^n(b_N, t))^2 - (\rho_N(b_N, t))^2 \right) + 2\psi(t)v_N (\rho_N^n(b_N, t) - \rho_N(b_N, t)) \right) dt \leq \\ & \|2\psi(t) + v_N(\rho_N^n(b_N, t) + \rho_N(b_N, t))\|_\infty \cdot \|v_N(\rho_N^n(b_N, t) - \rho_N(b_N, t))\|_{L^1}. \end{aligned}$$

Tangent vectors for numerical optimization

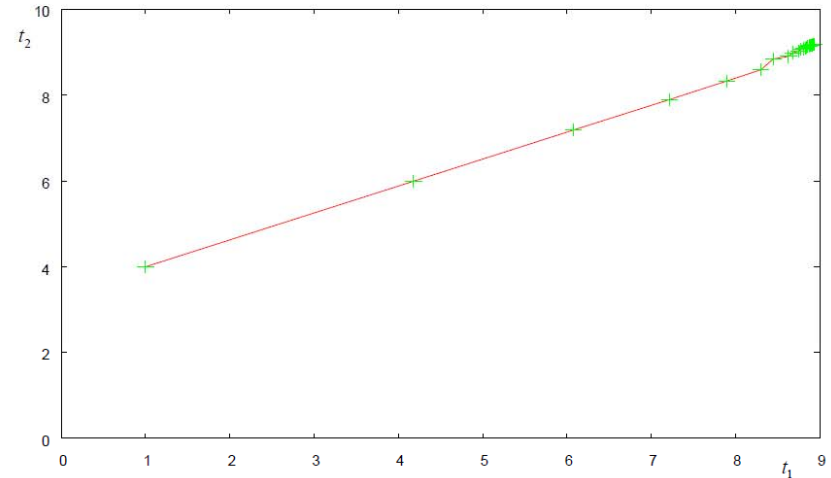
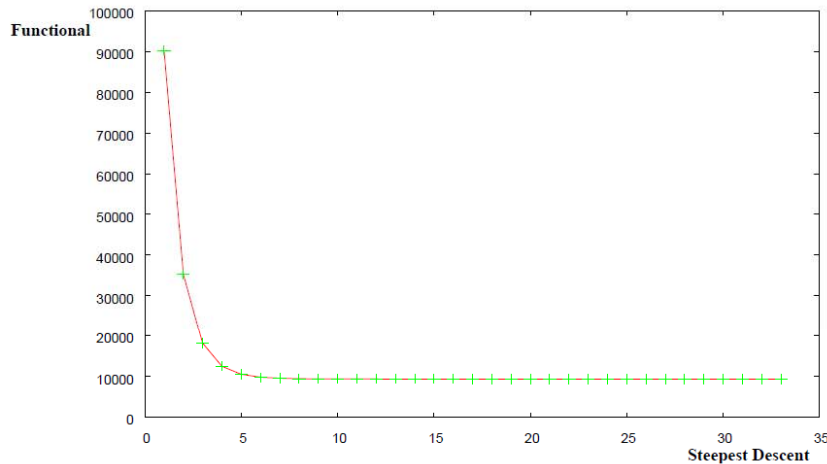
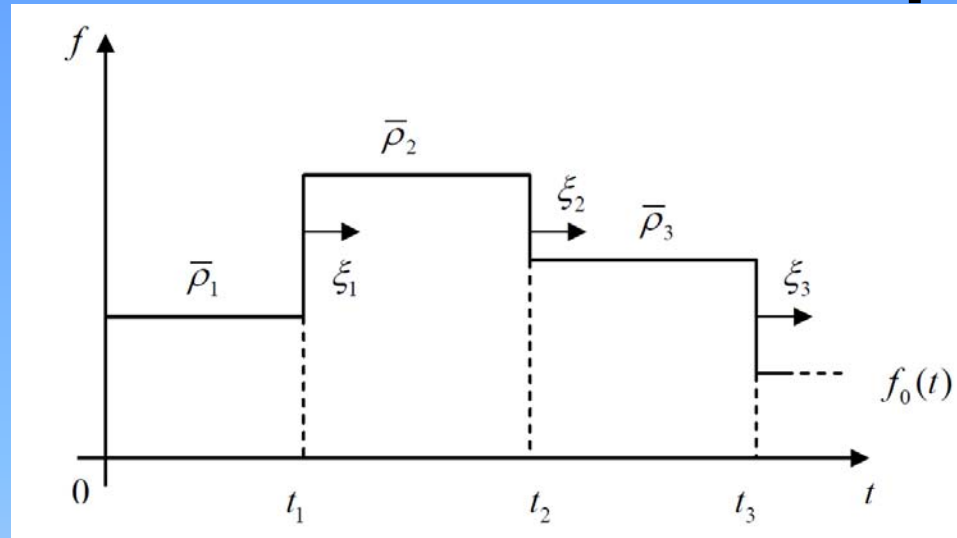


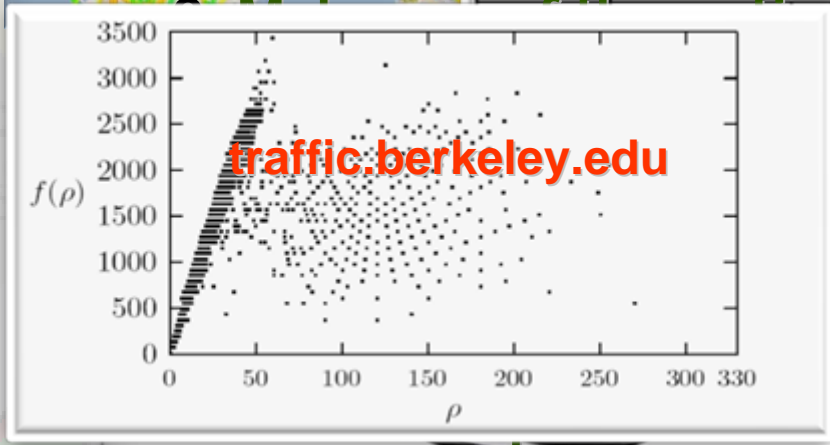
Figure 7: Supply chain with 11 arcs, case a . Left: J_1 versus iteration steps; right: "path" followed by the steepest descent algorithm in the plane (t_1, t_2) .

Cyber-infrastructure for information mobility

Network with 5000 roads parametrized by $[0,1]$,
 h space mesh size, T real time

1. Use simplified flux function with two characteristic speeds

CPU time				
h	G	FG	K3V	FSF
0.2	1.78 s	1.12 s	29.37 s	0.60 s
0.1	5.68 s	3.05 s	104.74 s	1.35 s
0.05	19.83 s	9.30 s	394.03 s	
0.025	73.86 s	31.40 s	1515.32 s	
$T = 30$				



3.32 s	85.71 s
0.12 s	309.55 s
7.93 s	1171.10 s
5.38 s	4527.80 s

FG = Fast Godunov
 etic, FSF = Fast
 ch zone σ



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CROWD DYNAMICS



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Alessia Marigo



Dirk Helbing



Seb Blandin



Gabriella Bretti



Michael



Rosanna Manzo



Axel Klar

SUPPLY CHAINS



Simone Goettl



Michael

ANIMAL GROUPS

Thank you for your attention!

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Real data

Problems :

1. Dimensionality: big networks
2. Data: measurements and elaboration



Manual counting



Satellite data

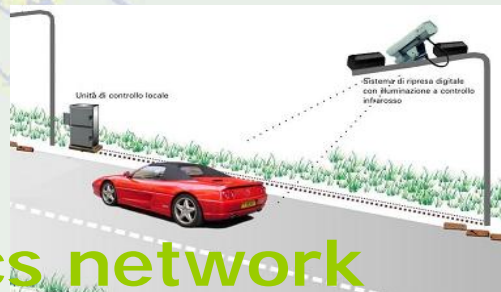


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1500 arcs network

Videocameras



Plates reading

NETWORK of SALERNO

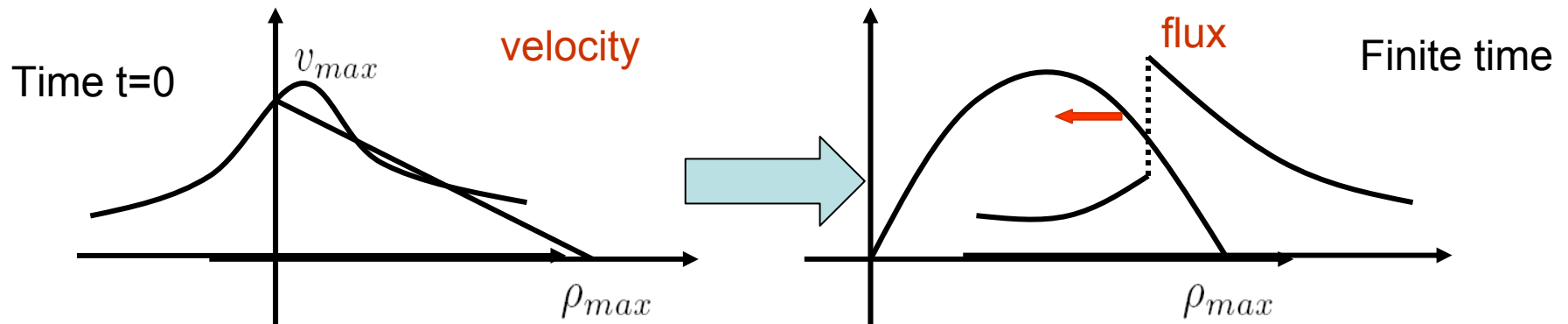
Lighthill-Whitham-Richards model

The flux is given by the density times the average velocity $f(t, x) = \rho(t, x) \cdot v(t, x)$

If we assume that the average velocity depends only on density $v(t, x) = v(\rho(t, x))$

Lighthill-Whitham-Richards model:

$$\rho_t + (\rho v(\rho))_x = 0, \quad \rho \in [0, \rho_{max}]$$

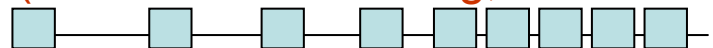


Non unique (weak) solutions

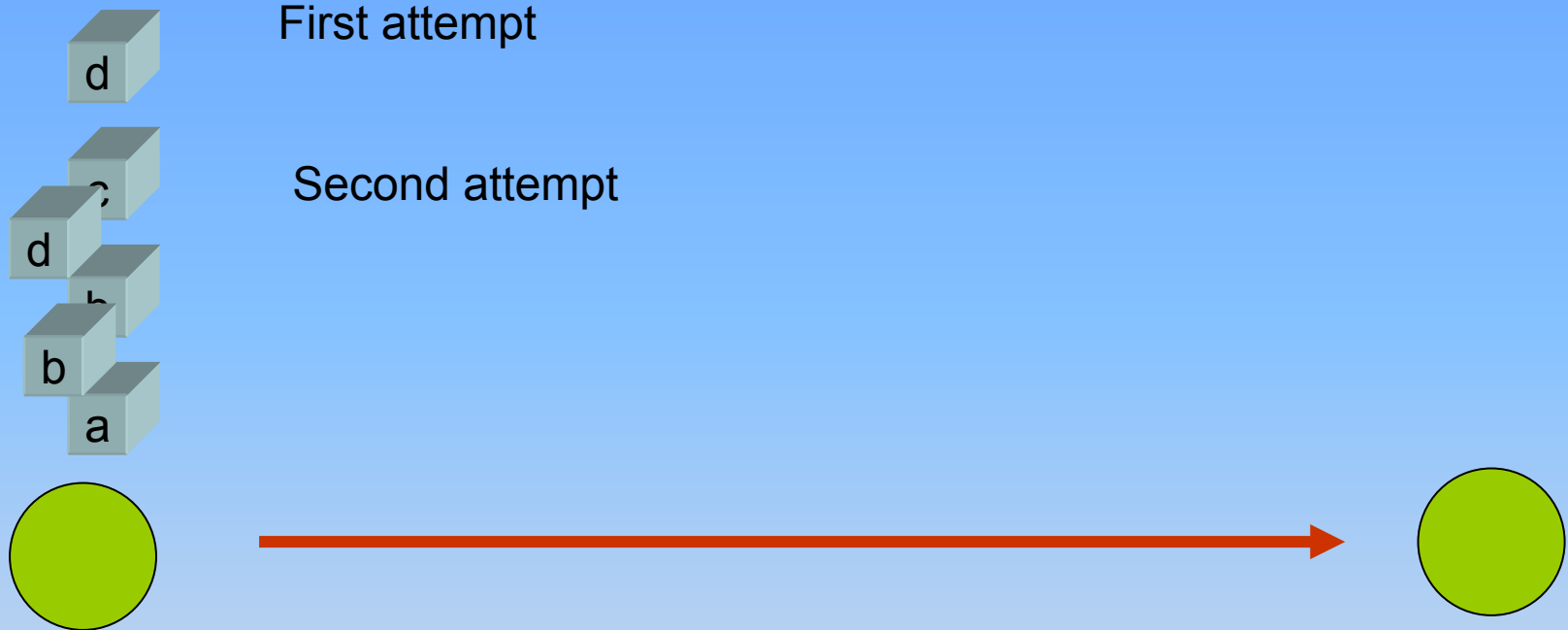


Entropy (gas dynamics)

(disorder is increasing, stable shocks)



Model for data networks



We essentially invert Rules (A) and (B), giving more importance to through flux than traffic distribution.

In the n th attempt $(1 - \mathcal{P}(R))\mathcal{P}(R)^{n-1}$ packets are sent and $\mathcal{P}(R)^n$ are lost.

Maximal fluxes on each line γ_i^{max} and γ_j^{max}

The through flux is $\Gamma = \min\{\sum_i \gamma_i^{max}, \sum_j \gamma_j^{max}\}$.